

Using Aco, Wsns Save Energy and Run for Longer

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ABSTRACT

WSNs are able to do both monitoring and communication. Due to their cheap cost and broad applicability, WSNs have attracted a lot of research. A WSN is a network of dispersed computing nodes that can perceive their surroundings by sending and receiving data in real time. There are several real-world monitoring uses for WSNs, including battlefield surveillance, environmental monitoring, health monitoring, and more. Power consumption and optimising network lifespan is the primary difficulty with WSNs. The study proposes an ACO-based strategy to WSNs that may increase the network's lifespan while reducing its energy usage. Ant Colony Optimization, or ACO, is a popular Meta heuristic that takes its cues from the real-world foraging strategies of ants. Stochastic constructive processes like the ant algorithm develop solutions by "walking" along a constructive network. In this research, we explore how to determine the largest possible number of interconnected covers across a variety of WSN deployments. Finding a single linked cover in a WSN has been the subject of many suggested strategies. Using the interconnected covers is a more direct approach to reducing energy usage and increasing the lifespan of the network. Several types of WSNs have benefited from the suggested method. In order to maximise network lifespan and reduce power consumption, the technique using LEACH and PARA, ACO was compared and found to be the most effective.

Keywords: Ant colony Optimization (ACO) algorithm, Energy efficiency, LEACH, Network lifetime, PARA, WSNs.

INTRODUCTION

One of the most challenging segments of the modern electronics industry is work on WSNs. It is predicted that these networks will be autonomous, low-power-challenging, context-aware, and adaptable [1]. Due to the potential for a large number of sensor nodes in a final application's environment [2], the deployment and maintenance of WSNs is a challenging issue. The purpose of a WSN is to connect users with useful information extracted from data collected by dispersed sensors in physical space. As a rule, sensor nodes rely on a small battery pack for their power needs. The ability to reduce energy consumption and increase the lifespan of WSNs is crucial to

their widespread use. One of the most important measures of a WSN's worth is its expected lifespan [3], which is how long it will take until the network stops meeting the needs of its applications. Due to the widespread use of nonrenewable batteries in WSN nodes, research aimed at extending the network's lifespan have emerged as one of the field's most pressing and difficult concerns. Densely placed nodes [4] in a WSN allow for a fraction of the nodes to solve the coverage and connection problems.

Because the remaining nodes can enter a sleep state to conserve energy, the lifetime of a WSN can be extended through careful planning of active intervals of nodes. However, at all times during the network's lifetime, the active nodes must form a connected cover to ensure adequate sensing coverage and network connectivity. In a WSN, the nodes are responsible for both monitoring and communication tasks. Nodes must provide enough sensing coverage of the target in order to perform the monitoring operation. The nodes must come together to create a network in order to gather and disseminate data through radio broadcasts, which is a requirement of the communication job. Because each linked cover must satisfy sensing coverage and network connection at the same time, determining the maximum number of connected covers is a challenging challenge. In this study, we focus on one class of WSNs and suggest a new method of activity planning to reduce energy consumption and extend the life of the network. In addition to area coverage, the method may also be used to situations involving discrete point coverage. Coverage is our main concern, with sensors and sinks making up the bulk of the nodes in the WSNs under consideration. The sensors keep an eye on the target and send their readings to the receiving nodes, or sinks. The monitoring data is sent to the target through the sinks. To this end, the following three requirements must be met by a linked cover in WSNs: 1) The item is completely covered by the sensors. All the data collected by the sensors is sent to the sinks, which are part of a wireless network, and are analysed there. The second limitation requires both sensors and sinks, therefore all three of these constraints are interrelated. Maximizing the number of sensor subsets under the coverage restriction and the number of sink subsets under the connectivity requirement are both more challenging than finding the greatest number of linked covers. While foraging through a construction graph, ACO's "ants" are really stochastic constructive methods that assemble solutions. The efficient search behaviour of ACO makes it a viable option for resolving combinatorial optimization challenges. To further expedite the search procedure, ACO makes use of search history and subject expertise. Many business and research issues have benefited from ACO algorithms. In order to enhance the power efficiency of unicasting, transmission, and data collection, this work suggested an ACO –based routing algorithm. Algorithms for increasing the lifespan of wireless sensor networks (WSNs) by identifying the largest number of linked covers have been developed, with an emphasis on routing challenges.

Work Cited

By fusing together communication technology, embedded computer, and sensor technology, Wireless Sensor Networks (WSNs) are a revolutionary kind of wireless communication network (WCN). Numerous sensor nodes are networked together over a wireless channel. Given that most sensor nodes rely on batteries for power, optimising the use of energy to maximise the network's lifespan has emerged as a fundamental challenge. Due to its importance in setting the stage for effective network communication, routing algorithms in WSNs have recently been a focus of

study. LEACH (Low- Energy Adaptive Clustering Hierarchy) is an application-level data distribution protocol that use clustering to extend the lifespan of a network. However, all of the CHs directly communicating with the base station might lead to the nodes dying out prematurely since the algorithm in the process of picking the CH does not account for the leftover energy of node. After proposing a fully distributed clustering (HEED) technique that takes into account the residual energy of each node to narrow down the pool of possible CHs, the authors of [5] suggest basing the final choice on the total cost of cluster communication. However, the communication cost is high since this technique needs several messages with the iteration in cluster radius. A decentralised routing method that aggregates data using the ANT algorithm is presented. The "ants" in this case are artificial agents designed to determine the best route to a given node, with the help of the positive feedback effect provided by the ANT algorithm. However, the technique is unable to resolve network energy load balancing. Through some tweaks to the Ant Colony Optimization algorithm, ACRA (ACO). Utilizing both the principal and secondary routes has led to a reduction in both energy usage and latency, but ant pathfinding is blind to factors other than pheromone, leading to congestion and a more inefficient use of resources overall. To improve its utility in WSNs, PARA, which incorporates the energy level and distance of transmission into the pheromone increment calculation. The method does not take into account the complete network's utilisation of the energy-stabilization problem, however. Given the proximity of the CH to the BTS, a large amount of data must be sent to

One solution to the issue is denoted by the set S , and if S is a subset of F , then S is a viable solution to the problem, as $F \supseteq C$ is the subset of feasible solutions. The goal is to locate a lowest cost feasible solution S^* , denoted by $S^* \in F$ and $Z(S^*) \leq Z(S), \forall S \in F$, where z is a cost function defined over the solution domain $Z: C \rightarrow R$. Colonies of computational ants travel from issue state to problem state, each stage representing a different solution to the problem. The way they go about is by using a flawed local choice policy based on two factors: trails and attraction. Each ant contributes to a solution by propelling it forward, but the whole colony's efforts are magnified by the speed with which they all work together. During the development stage, when an ant completes a solution, it evaluates the solution and updates the trail value of the components utilised to create the solution. The next ants' hunt will be guided by this pheromone report. Trail evaporation and daemon actions are two more processes seen in an ACO algorithm. To prevent an infinite accumulation of trails over any component, "trail evaporation" causes all trail values to depreciate with time. A daemon action may be used to order a local optimization procedure or update global information to be utilised in determining whether or not to bias the search process from a non-local standpoint. To be more precise, an ant is a straightforward computational agent that builds a workable answer to a problem by iterating over previous steps. States are considered to be partial answers to an issue. There is a cluster at the heart of the ACO algorithm, where each ant advances from state I to state j at each iteration, which corresponds to a more comprehensive partial solution. Each ant k at time t calculates a set $A(t)$ of possible expansions to its present state, and then probabilistically advances to one of them. These are the parameters for the probability distribution. The transition probability P_k for ant k from state I to state j is a function of the following two values: • the attractiveness η_{ij} of the move, as judged by some heuristic reflecting the a priori desire of that move;

We propose a new uneven clustering routing technique for Wireless Sensor Network based on the ant colony algorithm (ACO). This algorithm takes inspiration from both the ACA and the clustering routing algorithm already present in wireless sensor networks. The sensor nodes at the area level are divided using an uneven clustering technique such that the clusters closest to the sink have a lower size than those farther away from the sink. The selected CH then ran the combined clusters' data through the ACO algorithm to determine the best next steps. The algorithm's ability to stabilise the network's energy usage via simulation has proven successful in extending the network's lifespan.

ANT COLONY PERFORMANCE ANALYSIS

According to [6], a combinatorial optimization issue is one that is specified over a set $C = \{C_1, C_2, \dots, C_n\}$ of building blocks.

The degree to which an action has been successfully completed in the past is an indicator of how desirable it is.

When all ants have finished their answer, the trails are updated, with higher-level trails reflecting "excellent" solutions and lower-level trails reflecting "poor" ones.

Ant System

The original Ant System (AS) is notable because it served as the basis for many subsequent ant algorithms that, taken together, embody the ACO paradigm. For all infeasible moves (i.e., they are in the tabu list of ant k , that is a list containing all moves which are infeasible for ants k starting from state I the move probability distribution defines probabilities P_k to be equal to 0. Otherwise, they are computed by eq. (1), where α and β are user-defined parameters ($0 < \alpha, \beta < 1$): If (ij) is taboo, then (k) must be forbidden.

inspired by the desire to learn more about the algorithm and better understand its behaviour, which led to the discovery of several interesting results. In 1995, Gambardella and Dorigo presented Ant-Q, an extension of AS that incorporates certain principles from Q-learning. In 1996, they created Ant Colony System (ACS) [7], a reduced version of Ant-Q that retained about the same level of performance.

$P_{k,ij}$ efficiency, as evaluated by the number of sophisticated algorithms used and When updating the pheromone globally, however, ACS only uses the best solution found since the calculation began. While both ACS and AS use global updating to make a good route more appealing, the ACS process is superior since it avoids a lengthy convergence period by focusing the search precisely on the area around the greatest tour discovered so far. A local refresh of the pheromone applied during construction takes the place of the traditional final evaporation process in ACS.

To which E_i, E_j refers to the leftover energy at the i -th and j -th nodes, respectively. The ant system consists of a single main loop in which m ants build their solutions in simultaneously and then update the trail levels. Several parameters, including the value of, the ratio between the trail's importance and the attractiveness of the solution, the persistence of the trail, the starting level of

the trail (i.e., $\tau_{ij}(0)$), the number of ants (m), and the quality measure (Q) used to determine whether or not a solution is of high quality and low cost, all contribute to the algorithm's performance.

Methods of Organization in Ant Colonies

An effective population-based metaheuristic for approximating solutions to challenging optimization problems, ant colony optimization (ACO) was named after the swarm intelligence of the insects it mimics. With ACO, a colony of virtual "ants" works to find optimal solutions to a problem in optimization. Applying ACO involves recasting the optimization issue as one of finding the optimal route along a weighted graph. After starting with a blank graph, the artificial ants (henceforth ants) will gradually create solutions by traversing it. The Ant System algorithm was the first to take cues from the behaviour of actual ants. At first, AS was used to solve the travelling salesman issue, but it was no match for the best algorithms already in use for this task. However, he deserves credit for introducing ACO algorithms and demonstrating the promise of using artificial pheromone and artificial ants to propel the pursuit of perpetually superior solutions to intricate optimization challenges. In subsequent studies, belongs to the finest tour in the world, and the pheromone is still around. As a fascinating side effect of these local and global updating methods, the pheromone $\tau_{ij}(t)$ of each edge is inferior restricted by 0. The Max-Min-AS (MMAS) [8] is a similar method that proposes establishing boundaries on the value of the pheromone experiments.

Law for Changing States

When building a new solution, the state transition rule is used to determine the next step for each ant. In ACS, we offer a novel rule for transitioning between states termed pseudo-random-proportional.

In order to prevent the same node from being re-elected as CH, the pseudo-random proportional rule is utilised instead of the more extreme pseudo-random state choice process. We make the following changes to the $T(n)$ equation: in Ant System and the Q-learning-typically-used random-proportional-action-choice rule. If you follow the pseudo-random rule, you can be sure that the outcome will be the optimal one with probability q_0 (exploitation). The AS random-proportional rule employs a probability distribution over the future state that varies with τ_{ij} and η_{ij} . The pseudo-random-proportional state transition strategy used by ACS offers a straightforward method for striking a balance between probing uncharted territory and capitalising on established facts. highest quality

Clustering

After a node has through the prior steps to become a CH, its radius $c R$ is determined by the distance d determined during the cluster creation process. To write $c R$ as an equation, we have $d \leq c R$.

The probability with which a given state is selected is q_0 (with 0 indicating no selection and 1 indicating a random selection).

The maximum and lowest distances between network nodes and the base station, respectively, are

denoted by the notations d_{max} and d_{min} , while c is used to set the range of parameters (which may be anywhere from 0 to 1). The distance from node s_i to the DS is denoted by $d(s_i, DS)$, while the distance from the 0th node to the base s

Enhancing Efficiency with Hybridization

Significant symmetric and asymmetric TSP/ATSP instances were solved using ACS. An enhanced data structure called a candidate list is used for this in ACS. For each given city I a candidate list of length cl would include the cl most desired cities to visit. When deciding which city to relocate to, an ant in ACS first consults a candidate list informed by the state transition laws. In the event that the ant is unable to visit any of the cities on the shortlist, it will choose the closest accessible city only based on the heuristic value ij . Local optimization heuristic (hybridization) has been incorporated into ACS for TSP/ATSP to improve it. The idea is that whenever the ant generates a solution, that solution is then optimised to the local minimum using a local optimization heuristic based on an edge exchange strategy, such as 2-opt, 3-opt, or Lin-Kernighan. Pheromone trails are universally updated to reflect the new, optimal solutions, which are treated as the final solutions provided by the ants in this iteration. In order to compete with the state-of-the-art algorithm for the solution of TSP/ATSP problems, this ACS implementation combined a new pheromone management policy with a novel state transition approach and local search processes.

Clustering

The primary goal of the cluster route stage is information transfer. It is possible to perform multi-hop communication between clusters by utilising the ant colony optimization method to identify the ideal route between the CH, which will lower the energy consumption of the CH that is far from the base station. CH using the most efficient way for data transmission, therefore decreasing energy consumption and maximising the lifespan of the network.

cluster leadership election

First, the sensor node creates a random number between 0 and 1; if the generated number is less than the threshold $T(n)$, then the node in question is elected as the CH. If the current iteration's CH is the node in question, then set $T(n)$ to 0 at the beginning of the next iteration.

If there is no other candidate CH, a node in the cluster will identify itself as the CH and broadcast a message of competitive success to the other nodes in the cluster. If there are many possible CHs within a cluster, the residual energy of each node is calculated and compared; the node with the highest residual energy is selected as the CH and the good news is spread. Once the CH has been determined, the subordinate cluster may be determined based on the intensity of the received signal and the proper CH can be notified. Once the CH learns of the information, it employs the TDMA technique to schedule transmission times for the nodes, therefore completing the cluster.

Development of Algorithms

Place K ants at each CH; set the maximum number of hops (J_{Num}) to 1 and the maximum number of iterations (R_{max}) to C (a constant); initialise the pheromone matrix $Tabu$ to a matrix of size (n,n) ; record and store the generated path ($Tabu$); and record and store the best path (R_{best}) from CH to node (A_{city}).

Each ant looks for the next possible hop using a probability computation, inserts that node into A city, and then updates both A city and Tabu.

By moving from node I to node j, ants may alter the pheromone levels along the route.

Test each ant to see whether the iteration condition has been met. The search procedure continues until the ant satisfies the requirements, at which point the current optimum solution is saved.

SIMULATION RESULTS

An Overview of the Energy Model and Its Settings

The wireless transmitter module may adjust the amount of power it sends out based on the distance between the nodes, using the same energy model that we use in this research. Using MATLAB, we built an event-driven simulator. The precise nature of the simulation setting: The network consists of 200 sensor nodes, all of which are randomly dispersed in a 200m by 200m square. PacketLength =4000, ctrPacketLength =100, Rs =20 m, EDA =0.5 nJ/bit, Eelec =50 nJ/bit, fs =10 pJ/(bit•m²), mp =0.0013 pJ/(bit•m⁴), d0 =75 m, c =0.5, E0 =0.5J, R0 =30 m, =2, Analyzing Simulation Results You can see how our approach stacks up against LEACH and PAPA in terms of average energy usage per node across several simulation rounds. At 1150 rounds in LEACH, all 200 nodes are destroyed, using up 100J of energy; at the same time, PARA uses up 80J of energy; and our approach uses up 55J of energy, reducing energy use by 45J throughout the network.

Indicated in Graph 1 Is the Typical Energy Use The simulation results for the round-by-round survival rate of the nodes. While PARA's node survival rate was much lower than our algorithm's, it was significantly higher than LEACH's. The reason why this occurs is because LEACH does not account the remaining energy of the nodes when picking the CH and CH directly interact with base station, and then overly consume the energy of the CH. Although PARA employed the ant colony optimization technique to transmit data in multi-hop, it did not address the balanced energy consumption of the overall network. So that the CH near the base station will send enormous quantity of data, and will be early death.

CONCLUSION

In this research we have demonstrated a homogeneous clustering based ACO method for wireless sensor network that saves energy and maximise network lifespan. Ensuring a uniform node distribution across clusters extends the network's operational lifespan. A new cluster head is picked on the basis of the remaining energy of current cluster heads, holdback value, and closest hop distance of the node. The homogeneous method makes sure that every node is either a cluster leader or a member of one of the clusters in the wireless sensor network. The suggested clustering ACO method evenly distributes cluster members, which increases the network's longevity. What's more, under the proposed protocol, only cluster leaders, as opposed to all nodes, broadcast the cluster creation message. As a result, it extends the usefulness of sensor networks. Due to the limited battery life of sensors, achieving high energy efficiency is a significant problem when developing protocols for WSNs. The protocol was created with the intention of maximising the lifespan of the network by maximising the sensors' availability. Research gaps exist about the elements that influence cluster formation and CH communication. This study has significant

implications for the development of low-power wireless sensor networks. We employ an ant colony optimization technique to find the most efficient way for the CH to transfer data. According to the simulation findings, our method outperforms LEACH and PARA in terms of average energy usage, survival rate, and network life cycle duration.

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